

MATHEMATICS OLYMPIAD TRAINING SESSIONS

2023 Senior Mathematics Contest Problems with Some Solutions

1. For a string of A s and B s, define the substring replacement operation: $AB \rightarrow BBAA$. Starting from any string of A s and B s, is it always possible to perform a sequence of such replacement operations to obtain a string where all B s are to the left of all the A s.

Solution. Answer: No, it is not always possible after a sequence of replacement operations to obtain a string where all B s are to left of all the A s.

If the initial string contains a substring AAB or ABB , the string always contains a substring AAB or ABB , in which case there is always at least one A to the left of a B .

Proof. We suppose at some stage the string contains a substring AAB or ABB .

Case 1: The string before application of the replacement operation contains AAB .

After the replacement operation, either

AAB is unchanged; then: final string contains AAB ,

or

$AAB \rightarrow ABBAA$; then: final string contains ABB .

Case 2: The string before application of the replacement operation contains ABB .

After the replacement operation, either

ABB is unchanged. Then: final string contains ABB ,

or

$ABB \rightarrow BBAAB$. Then: final string contains AAB .

So, in both cases, after each replacement operation, the string contains AAB or ABB . \square

2. A positive integer m is given to Alice and Bob. Alice and Bob play a game:

Alice goes first, writing a non-zero digit on the board.

Then Bob and Alice alternate,

appending a digit to the front or back of the current number on the board, except that a digit appended at the front of the number must be non-zero.

Bob wins if at any time the number on the board is divisible by m .

- (i) Find the least m such that Alice can prevent Bob from winning.
- (ii) Same problem as (i), except Alice can write any $n \in \mathbb{N}$ at the start.

Solution. Answers: (i) $m = 12$. (ii) $m = 11$.

Proof. We either explain why, for a given m , Alice cannot prevent Bob from winning, or provide Alice's prevention strategy.

Case 1: $m \leq 10$ (for both (i) and (ii)).

After Alice writes starting number X , the possible numbers Bob can form:

$$\overline{Xd}$$

by appending $d \in \{0, 1, \dots, 9\}$, are 10 consecutive integer possibilities, (at least) one of which is divisible by $m \in \{1, 2, \dots, 10\}$.

Bob makes the appropriate choice and wins.

Case 2: $m = 11$.

For (i): After Alice writes $d \in \{1, 2, \dots, 9\}$,

Bob appends d obtaining: \overline{dd} which is divisible by 11 and wins.

For (ii): Here Alice can prevent Bob from winning.

Alice always leaves a number $X \equiv -1 \pmod{11}$,

with an odd number of digits, initially choosing 120.

If Bob responds by appending $d \neq 0$ at front of X ,

where X has k digits then

$$\begin{aligned} \overline{dX} &\equiv d \cdot 10^k + -1 \pmod{11} \\ &\equiv d \cdot (-1)^k + -1 \pmod{11} \\ &\equiv (d+1) \cdot -1 \pmod{11}, \text{ where } d+1 \in \{2, \dots, 10\} \\ &\not\equiv 0 \pmod{11} \end{aligned}$$

Alice responds by also appending $d \neq 0$ at front of X , so that

$$\overline{ddX} \equiv -1 \pmod{11}, \text{ since } 11 \mid \overline{dd} \cdot 10^k,$$

and \overline{ddX} again has an odd number of digits.

Otherwise, if Bob responds by appending $d \in \{0, 1, \dots, 9\}$ at back of X then

$$\begin{aligned} \overline{Xd} &\equiv 10 \cdot -1 + d \pmod{11} \\ &\equiv 1 + d \pmod{11}, \text{ where } d+1 \in \{1, 2, \dots, 10\} \\ &\not\equiv 0 \pmod{11} \end{aligned}$$

Alice responds by also appending d at back of X , so that

$$\begin{aligned} \overline{Xdd} &\equiv 100 \cdot -1 + d \cdot 11 \pmod{11} \\ &\equiv 1 \cdot -1 + 0 \pmod{11} \\ &\equiv -1 \pmod{11} \end{aligned}$$

and \overline{Xdd} again has an odd number of digits.

Case 3: $m = 12$ (for (i)).

Here Alice seeing X , can prevent Bob from winning,

by appending an appropriate $d \in \{1, 3, 5\}$ to the back of X .

Now, divisibility by $12 = 2^2 \cdot 3$, is equivalent to divisibility by both 4 and 3.

Firstly, since Alice forms an odd number \overline{Xd} ,

Bob must append an even digit c after d , (since $n : 12 \implies n : 2$).
But for \overline{Xdc} to be divisible by 4, we show Bob must append $c \in \{2, 6\}$:

$$\begin{aligned}\overline{Xdc} &\equiv \overline{dc} \pmod{4}, \text{ since } X \cdot 100 : 4 \\ &\equiv \overline{1c} \pmod{4}, \text{ since } 20 : 4 \implies 50 \equiv 30 \equiv 10 \pmod{4} \\ &\equiv 0 \pmod{4} \iff c \in \{2, 6\}.\end{aligned}$$

The choice d that Alice makes of 1, 3, 5 is to prevent Bob being able to make \overline{Xdc} divisible by 3:

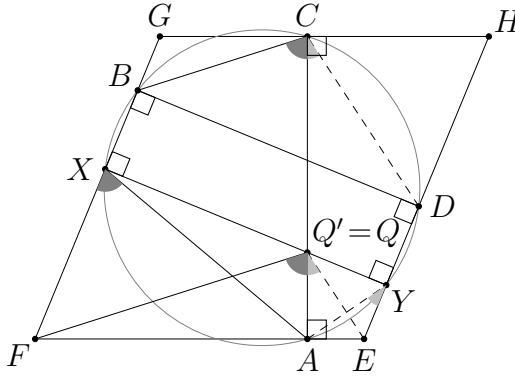
$X \pmod{3}$	Alice appends d	Bob appends c	$\overline{Xdc} \pmod{3}$
0	5	2, 6	1, 2
1	1	2, 6	1, 2
2	3	2, 6	1, 2

Initially X is empty, i.e. $X \equiv 0 \pmod{3}$, so that initially Alice chooses 5.

In any case, we see that Bob is prevented on every turn from winning. \square

3. Points A, B, C, D lie on sides EF, FG, GH, HE , respectively, of a parallelogram $EFGH$. Also $AC \perp EF, BD \perp FG, ABCD$ is cyclic, and Q is the point on AC such that $FQ \parallel BC$. Prove that $EQ \parallel DC$.

Proof 1. We will define Q' on AC , show $FQ' \parallel BC$, so that $Q' = Q$.
Then similarly show $EQ' \parallel DC$. Let



$$\begin{aligned}K &= \text{circumcircle}(ABCD) \\ \{X\} &= GF \cap K \setminus \{B\} \\ \{Y\} &= EH \cap K \setminus \{D\} \\ Q' &= XY \cap AC.\end{aligned}$$

Then

$DBXY$ is cyclic (circumcircle is K)
with adj. $\angle D = \angle B = 90^\circ$

$\therefore DBXY$ is a rectangle

$AFXQ'$ is cyclic (opp. $\angle A = \angle X (= 90^\circ)$ are suppl.)

$AXBC$ is cyclic (circumcircle is K)

$\therefore \angle FQ'A = \angle FXA$, angles on \widehat{FA} in cyclic $AFXQ'$
 $= \angle BCA$, int. opp. angle of ext. angle
of cyclic $AXBC$

$\therefore FQ' \parallel BC$, (corr. angles)

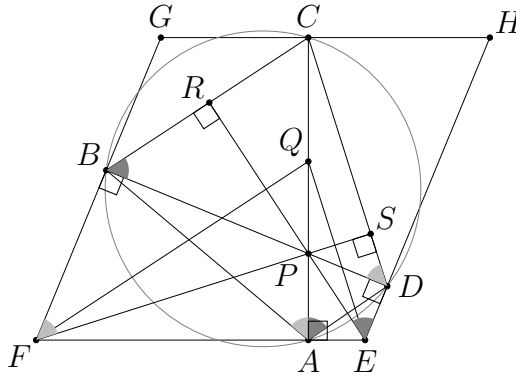
$\therefore Q' = Q$, since Q' on AC

Similarly, (map: $F \rightarrow E, X \rightarrow Y, B \rightarrow D$)

$\angle EQ'A = \angle EYA$, angles on \widehat{EA} in cyclic $AEQ'Y$
 $= \angle DCA$, int. opp. angle of ext. angle
of cyclic $AYDC$

$\therefore EQ = EQ' \parallel DC$, (corr. angles) \square

Proof 2. This time we show $FQ \parallel BC$ implies $P = \text{orthocentre}(EFQ)$ implies $EQ \parallel DC$.
Let



$K = \text{circumcircle}(ABCD)$
 $\ell = \text{line through } F, \parallel BC$
 $Q = \ell \cap CA$
 $P = BD \cap CA$
 $R = EP \cap BC.$
 $S = FP \cap DC.$

Then

$AEDP$ is cyclic (opp. $\angle A = \angle D (= 90^\circ)$ are suppl.)

$\therefore \angle DBR = \angle DBC$, same angle
 $= \angle DAC$, angles on \widehat{DC} in cyclic $ABCD$
 $= \angle DAP$, same angle
 $= \angle DEP$, angles on \widehat{DP} in cyclic $AEDP$
 $= \angle DER$, same angle

$\therefore BRDE$ is cyclic

$\therefore \angle BRE = \angle BDE$, angles on \widehat{BE} in cyclic $BRDE$
 $= 90^\circ$

Similarly, (map: $B \leftrightarrow D, E \rightarrow F, R \rightarrow S$)

$\angle BDS = \angle BFS$, $DSBE$ cyclic, $\angle DSF = \angle DBF = 90^\circ$

Now, $EP = ER \perp BC \parallel FQ$

$\therefore EP, AP$ are altitudes of $\triangle EFQ$

$\therefore P = \text{orthocentre}(EFQ)$

$\therefore FP$ is an altitude of $\triangle EFQ$

$\therefore FP \perp EQ$ and $FP \perp DC$

$\therefore EQ \parallel DC$ □

4. Counters are placed, one at a time, in the unit squares of an $n \times n$ grid such that:

- (i) a counter can only be placed in an empty unit square,
- (ii) the first counter can be placed in any unit square,
- (iii) each subsequent counter can only be placed in a unit square S if

the sum of the number of counters already in the same row as S and
the number of counters already in the same column as S is odd.

For each $n \geq 2$, find the smallest possible number of empty unit squares remaining after a sequence of such counter placements.

5. Let a_1, a_2, \dots, a_n and b_1, b_2, \dots, b_n be (not necessarily positive) real numbers, where $n \geq 2$. Let K, L be the maximum and minimum, respectively, of b_1, b_2, \dots, b_n .

Prove that $\sum_{i < j} a_i a_j |b_i - b_j| \leq \frac{1}{2}(K - L)(a_1 + a_2 + \dots + a_n)^2$,

where $\sum_{i < j}$ denotes the summing over all pairs (i, j) such that $1 \leq i < j \leq n$.